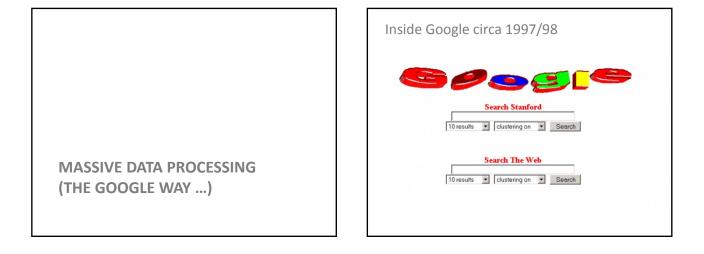
### CC5212-1 PROCESAMIENTO MASIVO DE DATOS OTOÑO 2015

Lecture 4: DFS & MapReduce I

Aidan Hogan aidhog@gmail.com Fundamentals of Distributed Systems

external sorts replication consistency consensus protocols cap theorem **availability** two phase commit fault tolerance distributed hash table partitions **client server** synchronous **client server** s



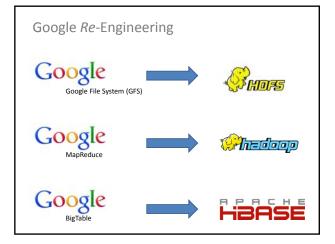
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# <section-header>Google Architecture (ca. 1998) Information Retrieval Crawling Inverted indexing Word-counts Ink-counts Ink-counts PageRank Updates

### **Google Engineering Google Engineering** Google File System • Massive amounts of data - Store data across multiple machines - Transparent Distributed File System • Each task needs communication protocols - Replication / Self-healing • Each task needs fault tolerance MapReduce • Multiple tasks running concurrently - Programming abstraction for distributed tasks – Handles fault tolerance Ad hoc solutions would repeat the same code - Supports many "simple" distributed tasks! • BigTable, Pregel, Percolator, Dremel ...



**GOOGLE FILE SYSTEM (GFS)** 

### What is a File-System?

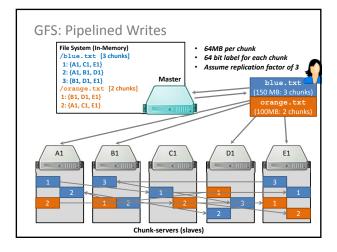
- Breaks files into chunks (or clusters)
- Remembers the sequence of clusters
- Records directory/file structure
- Tracks file meta-data
  - File size
  - Last access
  - Permissions
  - Locks

### What is a Distributed File-System

- Same thing, but distributed
- What would transparency / flexibility / reliability / performance / scalability mean for a distributed file system?
- Transparency: Like a normal file-system
- Flexibility: Can mount new machines
- Reliability: Has replication
- Performance: Fast storage/retrieval
- Scalability: Can store a lot of data / support a lot of machines

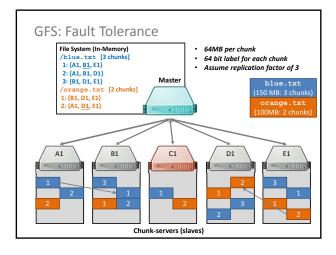
### Google File System (GFS)

- Files are huge
- Files often read or appended – Writes in the middle of a file not (really) supported
- Concurrency important
- Failures are frequent
- Streaming important



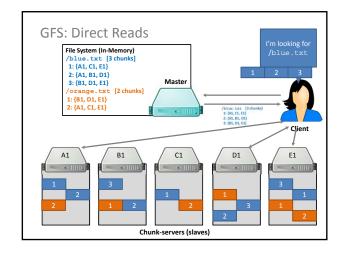
### GFS: Pipelined Writes (In Words)

- 1. Client asks Master to write a file
- 2. Master returns a primary chunkserver and secondary chunkservers
- 3. Client writes to primary chunkserver and tells it the secondary chunkservers
- 4. Primary chunkserver passes data onto secondary chunkserver, which passes on ...
- 5. When finished, message comes back through the pipeline that all chunkservers have written
  - Otherwise client sends again



### GFS: Fault Tolerance (In Words)

- Master sends regular "Heartbeat" pings
- · If a chunkserver doesn't respond
  - 1. Master finds out what chunks it had
  - 2. Master assigns new chunkserver for each chunk
  - 3. Master tells new chunkserver to copy from a specific existing chunkserver
- Chunks are prioritised by number of remaining replicas, then by demand

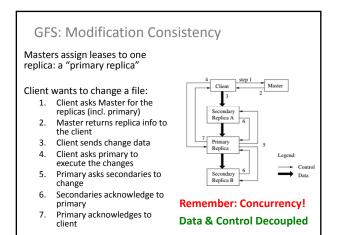


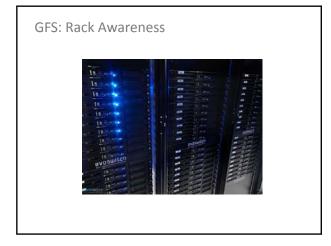
### GFS: Direct Reads (In Words)

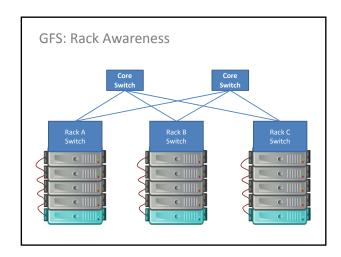
- 1. Client asks Master for file
- Master returns location of a chunk

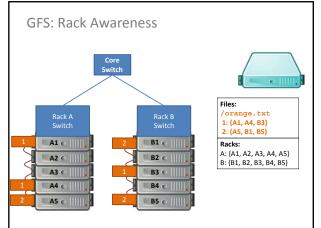
   Returns a ranked list of replicas
- 3. Client reads chunk directly from chunkserver
- 4. Client asks Master for next chunk

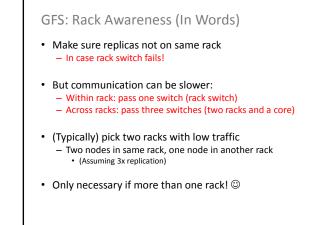
### Software makes transparent for client!











### **GFS:** Other Operations

Rebalancing: Spread storage out evenly

### **Deletion**:

- Just rename the file with hidden file name
  - To recover, rename back to original version
  - Otherwise, three days later will be wiped

Monitoring Stale Replicas: Dead slave reappears with old data: master keeps version info and will recycle old chunks

### GFS: Weaknesses?

# What do you see as the core weaknesses of the Google File System?

- Master node single point of failure

   Use hardware replication
   Logs and checkpoints!
- Master node is a bottleneck
  - Use more powerful machine
    Minimise master node traffic
- Master-node metadata kept in memory

   Each chunk needs 64 bytes
  - Chunk data can be queried from each slave

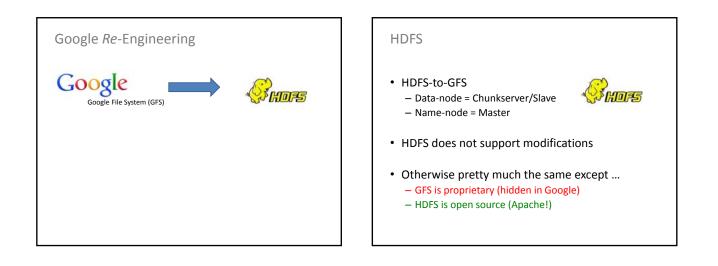
GFS: White-Paper

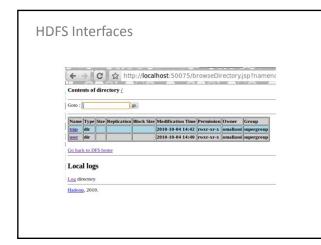
### The Google File System

Sanjay Ghemawat, Howard Gobioff, and Shun-Tak Leung Google\*

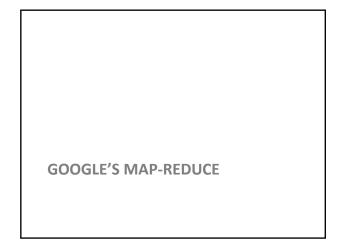
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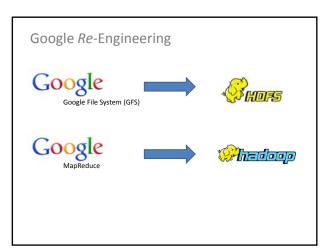
igh ag Whil  INTRODUCTION
 We have designed and implemented the Google Fi tem (GFS) to meet the rapidly growing demands of G data processing needs. GFS shares may of the sam as previous distributed file systems such as perfor scalability, reliability, and availability. However, its
 HADOOP DISTRIBUTED FILE SYSTEM (HDFS)





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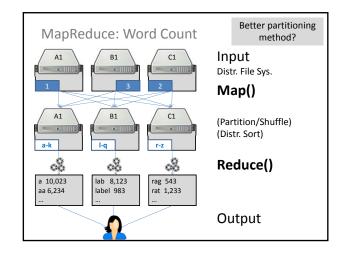


### MapReduce in Google

- Divide & Conquer:
- 1. Word count

How could we do a distributed top-k word count?

- 2. Total searches per user
- 3. PageRank
- 4. Inverted-indexing



### MapReduce (in more detail)

- 1. Input: Read from the cluster (e.g., a DFS)
  - Chunks raw data for mappers
  - Maps raw data to initial (key<sub>in</sub>, value<sub>in</sub>) pairs
     What might Input do in the word-count case?
- 2. <u>Map:</u> For each (key<sub>in</sub>, value<sub>in</sub>) pair, generate zero-to-many (key<sub>map</sub>, value<sub>map</sub>) pairs
  - key<sub>in</sub>/value<sub>in</sub> can be diff. type to key<sub>map</sub> /value<sub>map</sub>
     What might Map do in the word-count case?

### MapReduce (in more detail)

3. Partition: Assign sets of  $\mathrm{key}_{\mathrm{map}}$  values to reducer machines

How might Partition work in the word-count case?

- 4. Shuffle: Data are moved from mappers to reducers (e.g., using DFS)
- 5. Comparison/Sort: Each reducer sorts the data by key using a comparison function
  - Sort is taken care of by the framework

### MapReduce

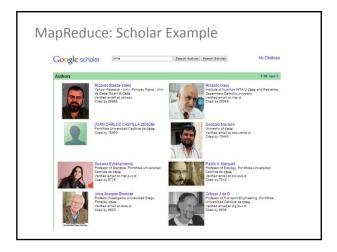
- - Typically zero-or-one outputs

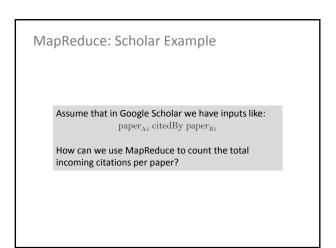
How might Reduce work in the word-count case?

7. Output: Merge-sorts the results from the reducers / writes to stable storage

### MapReduce: Word Count PseudoCode

```
function map(String name, String document):
    // name: document name
    // document: document contents
    for each word w in document:
        emit (w, 1)
function reduce(String word, Iterator partialCounts):
    // word: a word
    // partialCounts: a list of aggregated partial counts
    sum = 0
    for each pc in partialCounts:
        sum += ParseInt(pc)
    emit (word, sum)
```





### MapReduce as a Dist. Sys.

- Transparency: Abstracts physical machines
- Flexibility: Can mount new machines; can run a variety of types of jobs
- Reliability: Tasks are monitored by a master node using a heart-beat; dead jobs restart
- Performance: Depends on the application code but exploits parallelism!
- Scalability: Depends on the application code but can serve as the basis for massive data processing!

### MapReduce: Benefits for Programmers

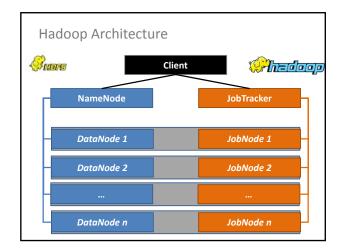
- Takes care of low-level implementation:
  - Easy to handle inputs and output
  - No need to handle network communication
  - No need to write sorts or joins
- Abstracts machines (transparency)
  - Fault tolerance (through heart-beats)
  - Abstracts physical locations
  - Add / remove machines
  - Load balancing

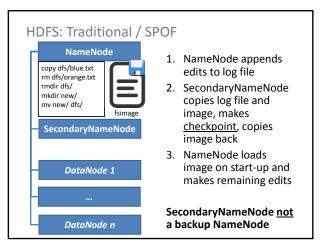
MapReduce: Benefits for Programmers

Time for more important things ...



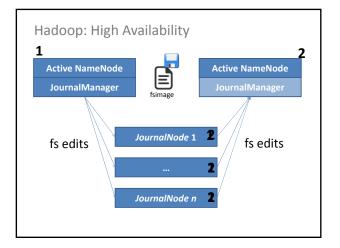
HADOOP OVERVIEW

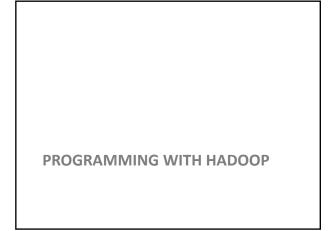


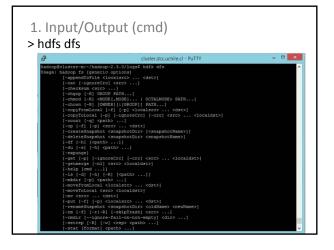


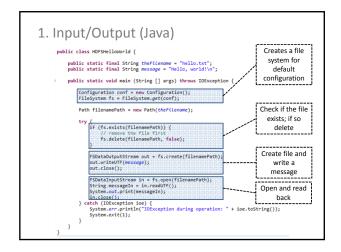
### What is the secondary name-node?

- Name-node quickly logs all file-system actions in a sequential (but messy) way
- Secondary name-node keeps the main fsimage file up-to-date based on logs
- When the primary name-node boots back up, it loads the fsimage file and applies the remaining log to it
- Hence secondary name-node helps make boot-ups faster, helps keep file system image up-to-date and takes load away from primary

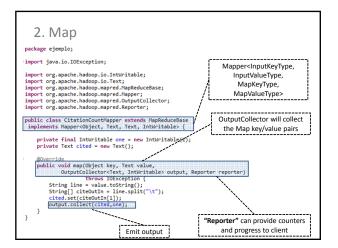


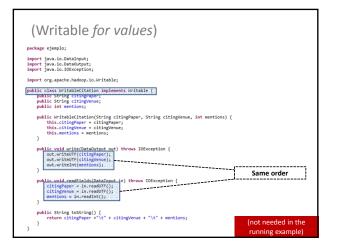


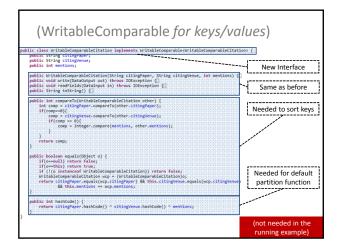


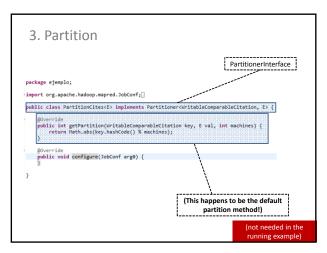


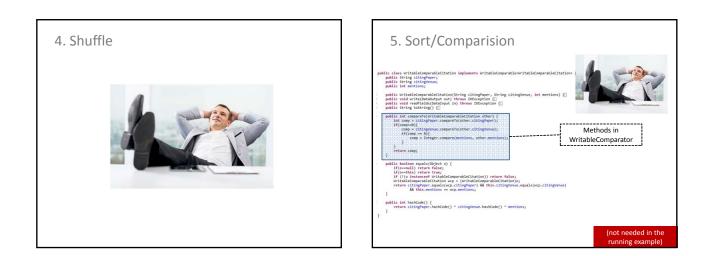
InputFormat:	Description:	Key:	Value:
TextInputFormat	Default format; reads lines of text files	The byte offset of the line	The line contents
KeyValueInputFormat	Parses lines into key, val pairs	Everything up to the first tab character	The remainder of the line
SequenceFileInputFormat	A Hadoop-specific high- performance binary format	user-defined	user-defined
SequenceFileInputFormat		user-defined	user-defin

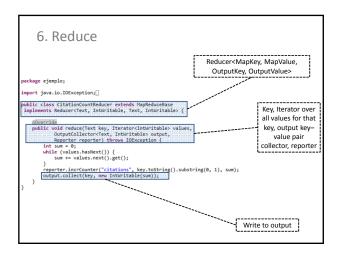


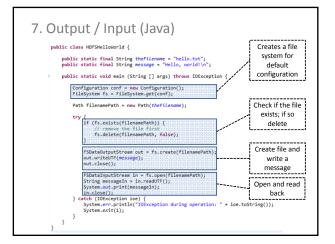




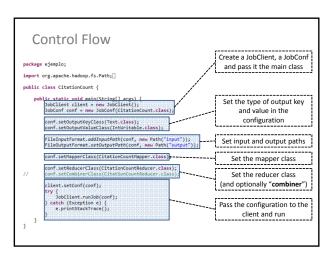








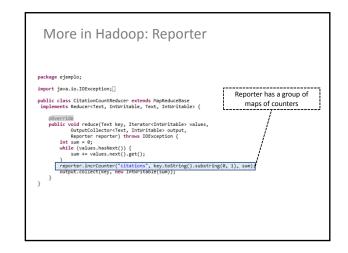
TextOutputFormat Default; writes lines in "key \t value" form	
SequenceFileOutputFormat Writes binary files suitable for reading into subseque	nt MapReduce jobs
NullOutputFormat Disregards its inputs	



### More in Hadoop: Combiner

- Map-side "mini-reduction"
- Keeps a fixed-size buffer in memory
- Reduce within that buffer

   e.g., count words in buffer
   Lessens bandwidth needs
- In Hadoop: can simply use Reducer class ③



More in Hadoop: Chaining Jobs

- · Sometimes we need to chain jobs
- In Hadoop, can pass a set of Jobs to the client
- x.addDependingJob(y)

### More in Hadoop: Distributed Cache

- Some tasks need "global knowledge"

   For example, a white-list of conference venues and journals that should be considered in the citation count
  - Typically small
- Use a distributed cache:
  - Makes data available locally to all nodes

# RECAP

### **Distributed File Systems**

- Google File System (GFS)
  - Master and Chunkslaves
  - Replicated pipelined writes
  - Direct reads
  - Minimising master traffic
  - Fault-tolerance: self-healing
  - Rack awareness
  - Consistency and modifications
- Hadoop Distributed File System
  - NameNode and DataNodes

# MapReduce

- 1. Input
- 2. Map
- 3. Partition
- 4. Shuffle
- 5. Comparison/Sort
  - 6. Reduce
  - 7. Output

# MapReduce/GFS Revision

- GFS: distributed file system – Implemented as HDFS
- MapReduce: distributed processing framework
  - Implemented as Hadoop

### Hadoop

- FileSystem
- Mapper<InputKey,InputValue,MapKey,MapValue>
- OutputCollector<OutputKey,OutputValue>
- Writable, WritableComparable<Key>
- Partitioner<KeyType,ValueType>
- Reducer<MapKey,MapValue,OutputKey,OutputValue>
- JobClient/JobConf
- ...